

# Primitive coherent configurations with very many automorphisms

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- 1 Background and motivation
- 2 Classification of PCCs with very many automorphisms
- 3 Structure theory for PCCs

# Coherent configurations

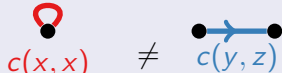
## Definition

Coherent configurations (CCs) **stable colorings under W-L**

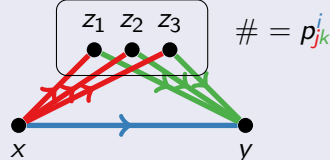
# Coherent configurations

## Definition

**CC**  $\mathfrak{X}$  on set  $\Omega$  is coloring  $c : \Omega \times \Omega \rightarrow \{\text{colors}\}$   
(edge-colored complete digraph with loops) s.t.

(1)   $c(x, x) \neq c(y, z)$

(2)   $c(x, y)$  determines  $c(y, x)$

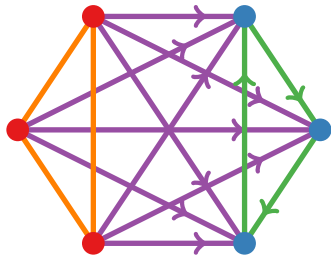
(3)   $\# = p_{jk}^i$

# Coherent configurations

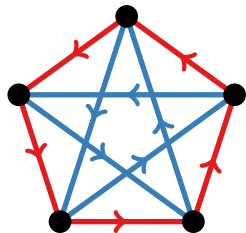
## Definition

CC is **homogeneous** if

(4) all vertices (loops) have same color



inhomogeneous



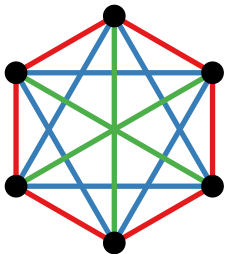
homogeneous

# Coherent configurations

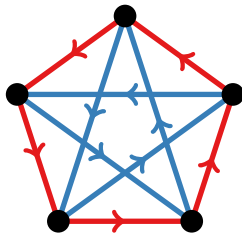
## Definition

Homogeneous CC  $\mathfrak{X}$  is **primitive (PCC)** if

(5) each constituent digraph  $\mathfrak{X}_i = (\Omega, c^{-1}(i))$  is connected

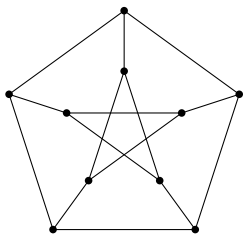


imprimitive



primitive

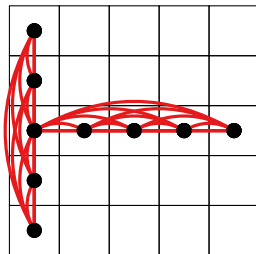
# Johnson and Hamming schemes



$J(m, k)$ :

Johnson scheme on  $\binom{[m]}{k}$

$$c(A, B) = |A \cap B|$$



$H(m, d)$ :

Hamming scheme on  $[m]^d$

$$c(w_1, w_2) = \text{Hamming distance}$$

Group action  $G \curvearrowright \Omega$  gives CC  $\mathfrak{X}(G)$  on  $\Omega$ :

$$c(u, v) = c(x, y) \text{ iff } (\exists g \in G)(u^g, v^g) = (x, y)$$

- $J(m, k) = \mathfrak{X}(S_m^{(k)})$
- $H(m, d) = \mathfrak{X}(S_m \wr S_d)$



Interested in **behavior of infinite families**  
as number of vertices  $n \rightarrow \infty$

E.g.,

- $|\text{Aut}(J(m, 2))| = \exp(\tilde{O}(n^{1/2}))$
- $|\text{Aut}(H(m, 2))| = \exp(\tilde{O}(n^{1/2}))$

# Symmetry vs. regularity

**Symmetry** - properties of automorphism group

**Regularity** - expressed in numerical parameters

All PCCs quite regular. . .

. . . but few PCCs are (very) highly symmetric!

## Conjecture (Babai)

$(\forall \varepsilon > 0)(\exists n_0)$

*every PCC with  $n \geq n_0$  verts. and  $\geq \exp(n^\varepsilon)$  automorphisms  
has primitive automorphism group*

## Theorem (Cameron 1981 via CFSG)

*Every primitive permutation group  $\Gamma$  has form*

$$(A_m^{(k)})^d \leq \Gamma \leq S_m^{(k)} \wr S_d$$

*or has order  $n^{O(\log n)}$  where  $n$  is degree*

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## Corollary (Under Babai's conjecture)

*Every PCC with  $\exp(n^\epsilon)$  automorphisms is Schurian  $\mathfrak{X}(\Gamma)$   
with  $\Gamma$  as above*

# PCCs with very many automorphisms

Theorem (Babai, 1981)

*PCCs other than trivial  $(K_n)$  have  $\leq \exp(\tilde{O}(n^{1/2}))$  automorphisms*

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Theorem (Sun–W 2015)

*PCCs other than trivial,  $J(m, 2)$ , and  $H(m, 2)$  have  $\leq \exp(\tilde{O}(n^{1/3}))$  automorphisms*

## Corollary (Sun–W 2015)

*Primitive permutation groups of degree  $n$   
except Cameron groups  
have order  $\leq \exp(\tilde{O}(n^{1/3}))$*

(CSFG-free proof of Cameron's theorem down to  $\exp(\tilde{O}(n^{1/3}))$ )



# Individualization/refinement: finding a combinatorial base

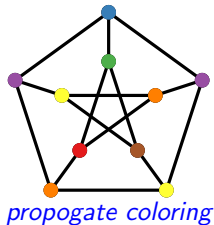
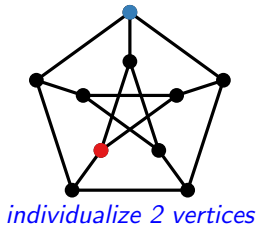
Set  $S$  is **base** if  $\text{Aut}(\mathfrak{X})_{(S)} = 1$ .

$\implies |\text{Aut}(\mathfrak{X})| \leq \exp(\tilde{O}(|S|))$

**Want**  $|S| \leq \tilde{O}(n^{1/3})$

Combinatorial relaxation:

- Fix vertices: assign unique colors (**individualize**)
- Estimate orbits: canonically **refine** coloring, examine color classes

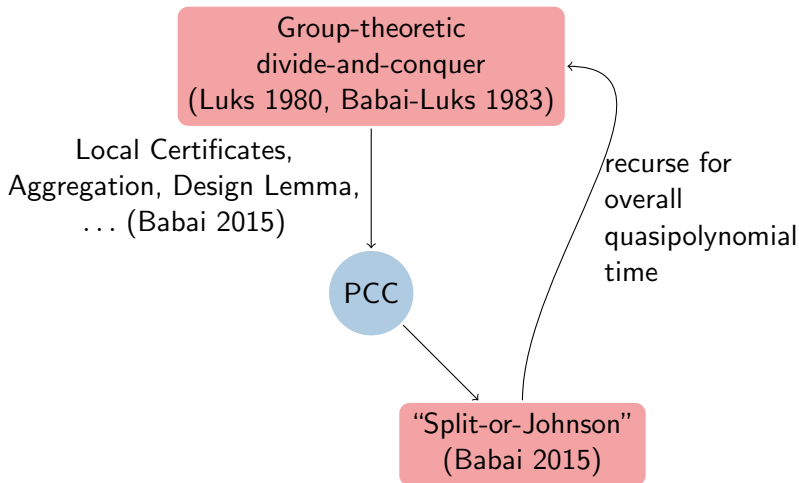


Goal: **discrete coloring** (all vertices end with unique colors)

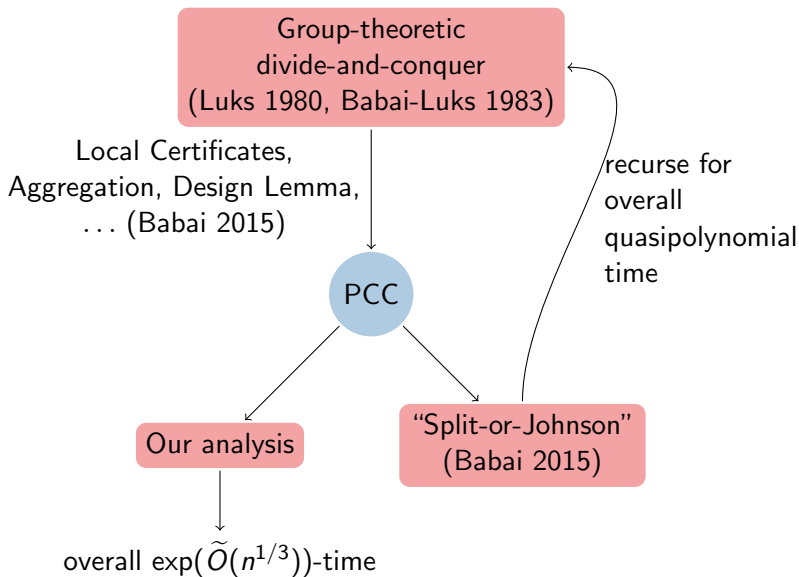
## Theorem (Sun–W 2015)

*Nontrivial PCCs other than  $J(m, 2)$  or  $H(m, 2)$  get discrete coloring from 1-WL after  $\tilde{O}(n^{1/3})$  individualizations*

# Babai's $\exp(\tilde{O}(n^{1/3}))$ Algorithm for General GI

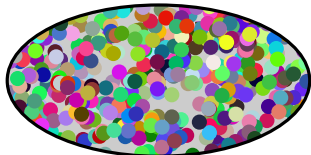


# Babai's $\exp(\tilde{O}(n^{1/3}))$ Algorithm for General GI



## Our analysis

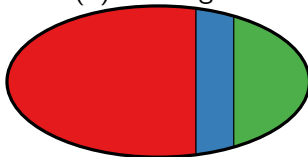
$J(m, 2)$  or  $H(m, 2)$ , or  
 $\tilde{O}(n^{1/3})$  indiv. gives:



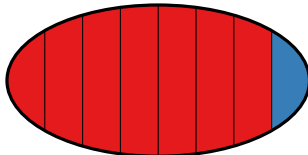
discrete coloring

## Babai's "Split-or-Johnson" Lemma

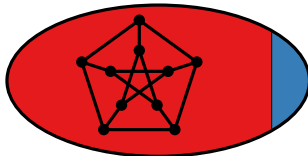
$\tilde{O}(1)$  indiv. gives:



no dominant color, or:



equipartition of dominant color, or:



$J(m, k)$  on dominant color

# Structure of primitive coherent configurations

## Theorem (Sun-W)

Let  $\mathfrak{X}$  be a PCC. One of these holds:

- 1 every edge-color has valency  $\leq n - n^{2/3}$  (no dominant color)
- 2 for some color  $i$ , we have  $\lambda_i = o(\sqrt{n})$
- 3 complement of dominant color has “clique geometry”

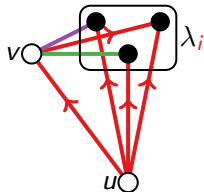
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$\lambda_i$  bound entails **vertex expansion**



# Structure of primitive coherent configurations

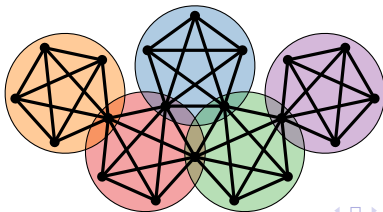
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## Definition

A **clique geometry** on graph  $G$  is set  $\mathcal{C}$  of maximal cliques in  $G$  s.t. every edge lies in unique clique





# Structure of primitive coherent configurations

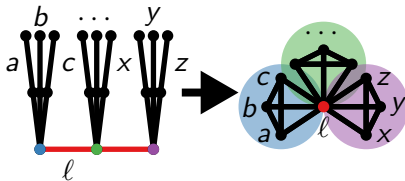
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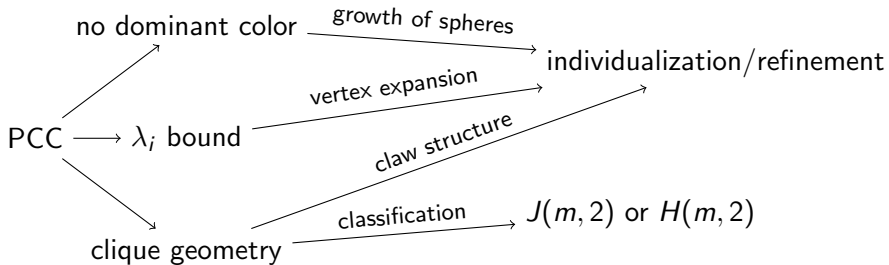
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Note:  $J(m, 2)$  and  $H(m, 2)$  have clique geometries

# Proof sketch

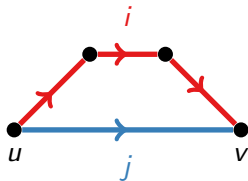


# When no color is overwhelmingly dominant

$\rho$  = co-valency of highest-valency color

**Assume**  $\rho \geq n^{2/3}$

$d = \max_{i,j} \text{dist}_i(j)$  (= distance in  $i$  between color- $j$  pair)



$$\text{dist}_i(j) = 3$$

Lemma (Babai 1981)

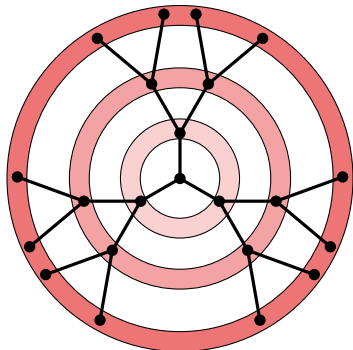
$O(nd \log n / \rho)$  individualizations suffice for discrete coloring

# When no color is overwhelmingly dominant

Lemma (Babai 1981)

$O(nd \log n/\rho)$  individualizations suffice for discrete coloring

We estimate growth of spheres



# When no color is overwhelmingly dominant

- $n_i =$  color  $i$  valency
- $\text{dist}_i(j) =$  distance in  $i$ th constituent between color- $j$  pair
- $S_i^{(\alpha)} =$  number of vertices in  $\alpha$ -sphere of  $i$ th constituent

## Lemma (Sun-W)

Let  $i, j$  off-diagonal colors and  $\delta = \text{dist}_i(j)$ . Then  $\forall 1 \leq \alpha \leq \delta - 2$

$$S_i^{(\alpha+1)} S_i^{(\delta-\alpha)} \geq n_i n_j$$

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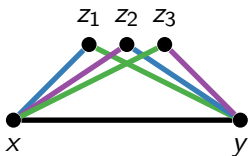
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**Enough when**  $\max n_i = \Omega(n)$

# When no color is overwhelmingly dominant

What if all  $n_i < \varepsilon n$ ? Analyze distinguishing number

$$D(u, v) = \#\{w : c(w, u) \neq c(w, v)\}$$



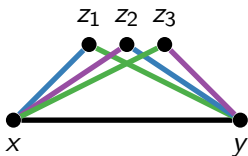
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Enough: all  $D(u, v)$  large

What if some  $D(u, v) = D(c(u, v))$  small?

# When no color is overwhelmingly dominant

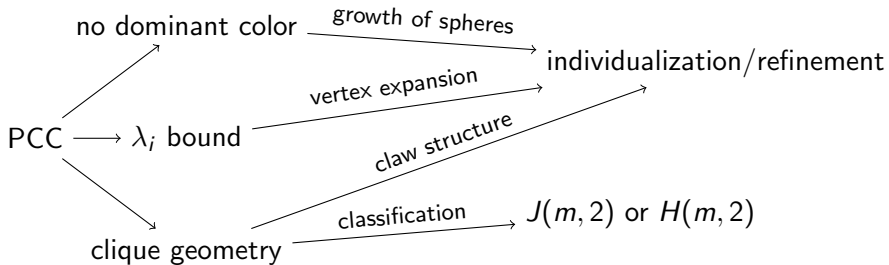
What if some  $D(u, v) = D(c(u, v))$  small?

Key tool

$D(i)$  small  $\implies$  “properties vary smoothly” along  $i$ -paths

In particular, valency and distinguishing number “vary smoothly”

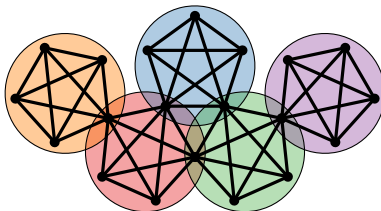
# Proof sketch



# Clique geometries in PCCs

## Definition

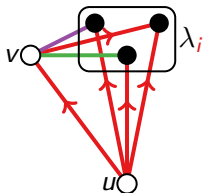
A **clique geometry** on graph  $G$  is set  $\mathcal{C}$  of maximal cliques in  $G$  s.t. every edge lies in unique clique



# Clique geometries in PCCs

Notation:

- $\mathfrak{X} = \text{PCC}$
- $\rho = \text{co-valency of highest-valency constituent}$
- $G = \text{complement graph of highest-valency constituent}$
- $\lambda_i = |\mathfrak{X}_i(u) \cap G(v)|$  where  $c(u, v) = i$



## Theorem (Sun-W)

$(\forall c > 0)(\exists \varepsilon > 0)$  s.t. if  $\lambda_i > cn^{1/2}$  and  $\rho < \varepsilon n^{2/3}$ ,  
then  $G$  has unique clique geometry.

Further, “asymptotically uniform”:

if  $c(u, v) = i$  in clique  $C$  then  $|G(u) \cap C| = \lambda_i + O(\rho\mu/\lambda_i)$

**Step 1:** Find “local clique partitions” in each color

## Definition

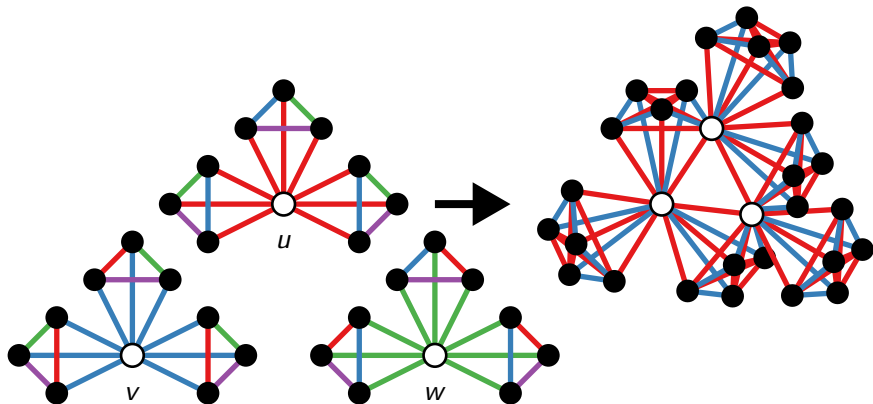
$\mathfrak{X}$  has  $I$ -local clique partitions for set  $I$  of colors if  $\forall u \in \Omega \exists \mathcal{P}$  partition of  $\mathfrak{X}_I(u)$  into maximal cliques of induced  $G$ -subgraph on  $\mathfrak{X}_I(u)$

Using assumptions on  $\rho$  and  $\lambda_i$ ,  $\mathfrak{X}$  has  $\{i\}$ -local clique partitions  $\forall i$  non-dominant by (Metsch 1991)

# Finding clique geometries

**Step 2:** Gradually strengthen clique partitions.

- 2a. Ensure cliques in  $\mathfrak{X}_I(u)$  maximal in  $G$  (not just induced subgraph)
- 2b. Ensure cliques agree at different vertices



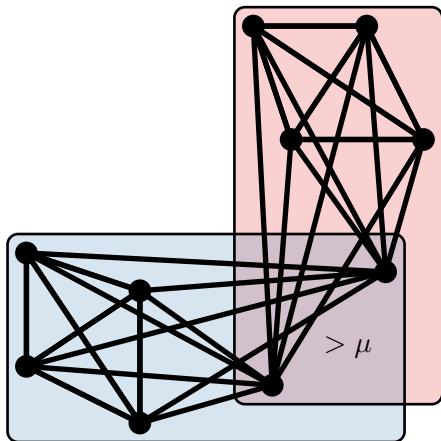


# Finding clique geometries

Let  $\mu = |G(u) \cap G(v)$  for  $u \neq v$   
in  $G$

## Observation

If cliques  $C_1, C_2$  in  $G$  have  
 $|C_1 \cap C_2| > \mu$  then  $C_1 = C_2$



**Step 3:** Put it together

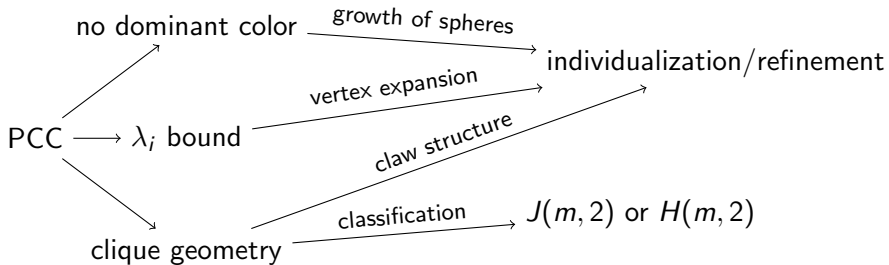
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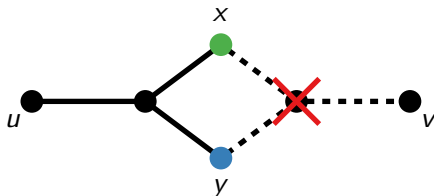
# Proof sketch



# Analysis of clique geometries

**Case 1:**  $\geq 3$  cliques at each vertex

Ubiquitous 3-claws in  $G$  give gadgets:



Individualizing  $x, y$  and refining distinguishes  $u$  and  $v$

**Case 2:**  $\leq 2$  cliques at each vertex

## Theorem (Sun-W)

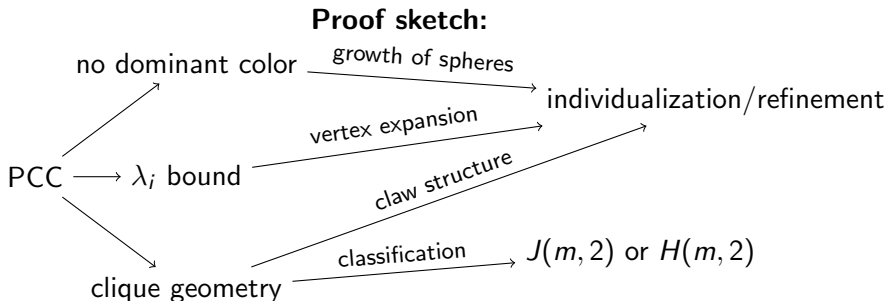
Suppose  $\rho < \varepsilon n^{2/3}$  and  $\exists$  asymptotically uniform clique geometry with  $\leq 2$  cliques at a vertex. Then  $X$  is one of these:

- 1 *trivial*
- 2 *Johnson  $J(m, 2)$*
- 3 *Hamming  $H(m, 2)$*
- 4 *rank 4 with a nonsymmetric color; symmetrization is  $H(m, 2)$ ;  $\leq n^O(\log n)$  automorphisms*

# Summary: Primitive Coherent Configurations

**Main result:** PCCs except trivial, Johnson, and Hamming have  $\leq \exp(\tilde{O}(n^{1/3}))$  automorphisms

*Previous best:*  $\exp(\tilde{O}(n^{1/2}))$ .



## Conjecture (Babai)

*Nontrivial PCCs other than Cameron schemes have  $\leq \exp(n^{o(1)})$  automorphisms*

### Easier cases?

- all valencies  $\leq \varepsilon n$
- exists non-symmetric color

(Definitely annoying cases)

Construct infinite family of PCCs with non-uniform clique geometries?